

PBES rewriters

Wieger Wesselink, Tim Willemse

January 19, 2019

This document describes several rewriters on PBES expressions. We assume that a data rewriter r is given that rewrites data expressions.

1 Simplifying rewriter

We define a simplifying PBES rewriter R recursively as follows. We assume that D is a non-empty data type, and we denote the free variables appearing in φ as $\text{free}(\varphi)$. We assume that a data rewriter r is given that rewrites data terms.

b	$\rightarrow r(b)$
$\neg\neg\varphi$	$\rightarrow \varphi$
$\varphi \wedge \text{true}$	$\rightarrow \varphi$
$\text{true} \wedge \varphi$	$\rightarrow \varphi$
$\varphi \wedge \text{false}$	$\rightarrow \text{false}$
$\text{false} \wedge \varphi$	$\rightarrow \text{false}$
$\varphi \wedge \varphi$	$\rightarrow \varphi$
$\varphi \vee \text{true}$	$\rightarrow \text{true}$
$\text{true} \vee \varphi$	$\rightarrow \text{true}$
$\varphi \vee \text{false}$	$\rightarrow \varphi$
$\text{false} \vee \varphi$	$\rightarrow \varphi$
$\varphi \vee \varphi$	$\rightarrow \varphi$
$\varphi \Rightarrow \psi$	$\rightarrow \neg\varphi \vee \psi$
$\forall_{d:D}.\varphi$	$\rightarrow \varphi$ if $d \notin \text{free}(\varphi)$
$\forall_{d:D}.\neg\varphi$	$\rightarrow \neg\exists_{d:D}.\varphi$
$\forall_{d:D}.\varphi \wedge \psi$	$\rightarrow \forall_{d:D}.\varphi \wedge \forall_{d:D}.\psi$
$\forall_{d:D}.\varphi \vee \psi$	$\rightarrow (\forall_{d:D}.\varphi) \vee \psi$ if $d \notin \text{free}(\psi)$
$\forall_{d:D}.\varphi \vee \psi$	$\rightarrow \varphi \vee (\forall_{d:D}.\psi)$ if $d \notin \text{free}(\varphi)$
$\exists_{d:D}.\varphi$	$\rightarrow \varphi$ if $d \notin \text{free}(\varphi)$
$\exists_{d:D}.\neg\varphi$	$\rightarrow \neg\forall_{d:D}.\varphi$
$\exists_{d:D}.\varphi \vee \psi$	$\rightarrow \exists_{d:D}.\varphi \vee \exists_{d:D}.\psi$
$\exists_{d:D}.\varphi \wedge \psi$	$\rightarrow (\exists_{d:D}.\varphi) \wedge \psi$ if $d \notin \text{free}(\psi)$
$\exists_{d:D}.\varphi \wedge \psi$	$\rightarrow \varphi \wedge (\exists_{d:D}.\psi)$ if $d \notin \text{free}(\varphi)$
$X(e)$	$\rightarrow X(r(e))$

where φ and ψ are arbitrary pbes expressions, b is a data term of data sort \mathbb{B} , true and false are elements of data sort \mathbb{B} , X is a predicate variable, e consists of zero or more data sorts and d is a data variable of sort D .

Simplify The pbes expression rewrite system SIMPLIFY [Luc Engelen, 2007] consists of the following rules¹:

$$\begin{aligned} false \wedge x &\rightarrow false \\ x \wedge false &\rightarrow false \\ true \wedge x &\rightarrow x \\ x \wedge true &\rightarrow x \\ \neg true &\rightarrow false \\ \neg false &\rightarrow true \\ ITE(true, x, y) &\rightarrow x \\ ITE(false, x, y) &\rightarrow y \\ x = x &\rightarrow true \\ y = x &\rightarrow x = y, \text{ provided } y \succ x \end{aligned}$$

¹Todo: reformulate this rewrite system.

2 PFNF Rewriter

Definition 1 A predicate formula is said to be in Predicate Formula Normal Form (PFNF) if it has the following form:

$$\mathbf{Q}_1 v_1:V_1 \cdots \mathbf{Q}_n v_n:V_n. h \wedge \bigwedge_{i \in I} \left(g_i \implies \bigvee_{j \in J_i} X^j(e^j) \right)$$

where $X^j \in \chi$ (χ is a countable of sorted predicate variables), $\mathbf{Q}_i \in \{\forall, \exists\}$, I is a (possibly empty) finite index set, each J_i is a non-empty finite index set, and h and every g_i are simple formulae.

Note that here J_i is used to index a set of occurrences of not necessarily different variables. For instance, $(n > 0 \implies X(3) \vee X(5) \vee Y(6))$ is a formula complying to the definition of PFNF. So long as it does not lead to confusion, we stick to the convention to drop the typing of the quantified variables v_i . An algorithm to compute a PFNF is:

$$\begin{aligned} p(c) &=_{def} c \\ p(X(d)) &=_{def} X(d) \\ p(\forall x:D.\varphi) &=_{def} \forall x:D.p(\varphi) \\ p(\exists x:D.\varphi) &=_{def} \exists x:D.p(\varphi) \\ \\ p(\varphi \wedge \psi) &=_{def} \mathbf{Q}_1^\varphi \cdots \mathbf{Q}_{n^\varphi}^\varphi \mathbf{Q}_1^\psi \cdots \mathbf{Q}_{n^\psi}^\psi. (h^\varphi \wedge h^\psi) \\ &\quad \wedge \bigwedge_{i \in I^\varphi \cup I^\psi} \left(g_i \implies \bigvee_{j \in J_i} X^j(e^j) \right) \\ \\ p(\varphi \vee \psi) &=_{def} \mathbf{Q}_1^\varphi \cdots \mathbf{Q}_{n^\varphi}^\varphi \mathbf{Q}_1^\psi \cdots \mathbf{Q}_{n^\psi}^\psi. (h^\varphi \vee h^\psi) \\ &\quad \wedge \bigwedge_{i \in I^\varphi} \left(\neg h^\psi \wedge g_i \implies \bigvee_{j \in J_i} X^j(e^j) \right) \\ &\quad \wedge \bigwedge_{i \in I^\psi} \left(\neg h^\varphi \wedge g_i \implies \bigvee_{j \in J_i} X^j(e^j) \right) \\ &\quad \wedge \bigwedge_{i \in I^\varphi, k \in I^\psi} \left((g_i \wedge g_k) \implies \bigvee_{j \in J_i, m \in J_k} X^j(e^j) \vee X^m(e^m) \right) \end{aligned}$$

where

$$\begin{aligned} p(\varphi) &= \mathbf{Q}_1^\varphi \cdots \mathbf{Q}_{n^\varphi}^\varphi. h^\varphi \wedge \bigwedge_{i \in I^\varphi} \left(g_i \implies \bigvee_{j \in J_i} X^j(e^j) \right) \\ p(\psi) &= \mathbf{Q}_1^\psi \cdots \mathbf{Q}_{n^\psi}^\psi. h^\psi \wedge \bigwedge_{i \in I^\psi} \left(g_i \implies \bigvee_{j \in J_i} X^j(e^j) \right), \end{aligned}$$

under the assumption that I^φ and I^ψ are disjoint, and $v_i^\varphi \neq v_j^\psi$ for all i, j .

3 OnePointRule Quantifier Elimination Rewriter

The function Eq computes a set of equalities and inequalities for an expression φ , such that the following holds:

$$Eq(\varphi) = (\{b_1 = e_1, \dots, b_n = e_n\}, W) \Rightarrow \varphi = \psi \wedge \bigwedge_{i=1}^n (b_i = e_i)$$

$$Eq(\varphi) = (V, \{b_1 \neq e_1, \dots, b_n \neq e_n\}) \Rightarrow \varphi = \psi \vee \bigvee_{i=1}^n (b_i \neq e_i)$$

for some expression ψ .

The function Eq is inductively defined as follows:

$$\begin{aligned} Eq(b) &= (\{b = true\}, \emptyset) \\ Eq(d = e) &= \begin{cases} (\{d = e\}, \emptyset) & \text{if } d \notin FV(e) \\ (\emptyset, \emptyset) & \text{otherwise} \end{cases} \\ Eq(e = d) &= Eq(d = e) \\ Eq(d \neq e) &= \begin{cases} (\emptyset, \{d \neq e\}) & \text{if } d \notin FV(e) \\ (\emptyset, \emptyset) & \text{otherwise} \end{cases} \\ Eq(e \neq d) &= Eq(d \neq e) \\ Eq(\neg\varphi) &= swap(Eq(\varphi)) \\ Eq(\varphi \wedge \psi) &= join_and((Eq(\varphi), Eq(\psi))) \\ Eq(\varphi \vee \psi) &= join_or((Eq(\varphi), Eq(\psi))) \\ Eq(\varphi \Rightarrow \psi) &= join_or((swap(Eq(\varphi)), Eq(\psi))) \\ Eq(\forall x.\varphi) &= delete(x, Eq(\varphi)) \\ Eq(\exists x.\varphi) &= delete(x, Eq(\varphi)) \\ Eq(\varphi) &= (\emptyset, \emptyset) \text{ otherwise} \end{aligned}$$

where b is a boolean data variable, d is an arbitrary data variable, e is a boolean data expression, and

$$\begin{aligned} swap((V, W)) &= (W, V) \\ join_and((V_1, W_1), (V_2, W_2)) &= (V_1 \cup V_2, W_1 \cap W_2) \\ join_or((V_1, W_1), (V_2, W_2)) &= (V_1 \cap V_2, W_1 \cup W_2) \\ delete(x, (V, W)) &= \begin{cases} (V_1, W_1) \text{ where} \\ V_1 = \{d = e \in V \mid d \neq x \wedge x \notin FV(e)\} \\ W_1 = \{d \neq e \in W \mid d \neq x \wedge x \notin FV(e)\} \end{cases} \end{aligned}$$

We define the OnePointRule rewriter R inductively as follows:

$$\begin{aligned} R(\neg\varphi) &= \neg R(\varphi) \\ R(\varphi \wedge \psi) &= R(\varphi) \wedge R(\psi) \\ R(\varphi \vee \psi) &= R(\varphi) \vee R(\psi) \\ R(\varphi \Rightarrow \psi) &= R(\varphi) \Rightarrow R(\psi) \\ R(\forall x.\varphi) &= \begin{cases} R(\varphi)[x := e] & \text{if } x \neq e \in W, \text{ where } (V, W) = Eq(\varphi) \\ \forall x.R(\varphi) & \text{otherwise} \end{cases} \\ R(\exists x.\varphi) &= \begin{cases} R(\varphi)[x := e] & \text{if } x = e \in V, \text{ where } (V, W) = Eq(\varphi) \\ \exists x.R(\varphi) & \text{otherwise} \end{cases} \\ R(\varphi) &= \varphi \text{ otherwise} \end{aligned}$$

4 Quantifier Inside Rewriter

This rewriter was originally specified by Jan Friso Groote. We define a rewriter that pushes universal and existential quantifiers into an expression. A typical example is

$$\exists x.\forall y.(f(y, y) \wedge (f(x, y) \vee f(x, x)))$$

which is rewritten to

$$(\forall y.f(y, y)) \wedge ((\exists x.\forall y.f(x, y)) \vee \exists x.f(x, x)).$$

If quantifiers are pushed inside formulas as much as possible, quantifier elimination leads to smaller expressions and the one-point rewriter is more often applicable. There might be cases where the application of this push-quantifiers-inside-rewriter can also have averse effects.

Below the definition of the rewrite rules are given. It is assumed that it is cheap to obtain the free variable in each term, which are denoted by $vars(\phi)$ for a term ϕ . If not, care needs to be taken as in the formulation below calculating the variables of each term recursively on a by need basis can be very expensive.

Let ϕ be a pbes expression. We define $R(\phi)$ using the functions $R_{\forall}(V, \phi)$ and $R_{\exists}(V, \phi)$ where V is a set of typed variables. The definition of R employs the structure of a formula:

$$\begin{aligned} R(\neg\phi) &= \neg R(\phi) \\ R(\phi \wedge \psi) &= R(\phi) \wedge R(\psi) \\ R(\phi \vee \psi) &= R(\phi) \vee R(\psi) \\ R(\phi \Rightarrow \psi) &= R(\phi) \Rightarrow R(\psi) \\ R(\forall W.\phi) &= R_{\forall}(W, R(\phi)) \\ R(\exists W.\phi) &= R_{\exists}(W, R(\phi)) \\ R(\phi) &= \phi \text{ otherwise.} \end{aligned}$$

Here W is a set of typed variables.

The rewrite function R_{\forall} is defined by

$$\begin{aligned} R_{\forall}(V, \neg\phi) &= \neg R_{\exists}(V, \phi) \\ R_{\forall}(V, \phi \wedge \psi) &= R_{\forall}(V, \phi) \wedge R_{\forall}(V, \psi) \\ R_{\forall}(V, \bigvee_i \phi_i) &= \forall W. (R_{\forall}(V \setminus W, \bigvee_{i \in I_1} \phi_i) \vee \bigvee_{i \in I_2} R(\phi_i)) \\ &\quad \text{where } W = vars(\phi_j) \cap V \text{ for some } j \text{ such that } |W| \text{ is minimal,} \\ &\quad \text{and where } I_1 = \{i \mid vars(\phi_i) \cap V \not\subseteq W\}, I_2 = \{i \mid vars(\phi_i) \cap V \subseteq W\} \\ R_{\forall}(V, \phi \Rightarrow \psi) &= \forall W. (R_{\exists}(V \setminus W, \phi) \Rightarrow R_{\forall}(V \setminus W, \psi)), \text{ where } W = V \cap vars(\phi) \cap vars(\psi) \\ R_{\forall}(V, \forall W.\phi) &= R_{\forall}(V \cup W, \phi) \\ R_{\forall}(V, \phi) &= \forall V \cap vars(\phi). \phi \text{ otherwise.} \end{aligned}$$

The rewrite function R_{\exists} is defined by

$$\begin{aligned} R_{\exists}(V, \neg\phi) &= \neg R_{\forall}(V, \phi) \\ R_{\exists}(V, \bigwedge_i \phi_i) &= \exists W. (R_{\exists}(V \setminus W, \bigwedge_{i \in I_1} \phi_i) \wedge \bigwedge_{i \in I_2} R(\phi_i)) \\ &\quad \text{where } W = vars(\phi_j) \cap V \text{ for some } j \text{ such that } |W| \text{ is minimal,} \\ &\quad \text{and where } I_1 = \{i \mid vars(\phi_i) \cap V \not\subseteq W\}, I_2 = \{i \mid vars(\phi_i) \cap V \subseteq W\} \\ R_{\exists}(V, \phi \vee \psi) &= R_{\exists}(V, \phi) \vee R_{\exists}(V, \psi) \\ R_{\exists}(V, \phi \Rightarrow \psi) &= R_{\forall}(V, \phi) \Rightarrow R_{\exists}(V, \psi) \\ R_{\exists}(V, \exists W.\phi) &= R_{\exists}(V \cup W, \phi) \\ R_{\exists}(V, \phi) &= \exists V \cap vars(\phi). \phi \text{ otherwise.} \end{aligned}$$

5 Quantifier Expansion Rewriter

5.1 Conjunctions / disjunctions

The function CD computes a sequence of conjunctions and disjunctions for an expression φ , such that the following holds:

$$CD(\varphi) = ([\varphi_1, \dots, \varphi_n], W) \Rightarrow \varphi = \bigwedge_{i=1}^n \varphi_i$$

$$CD(\varphi) = (V, [\varphi_1, \dots, \varphi_n]) \Rightarrow \varphi = \bigvee_{i=1}^n \varphi_i.$$

The function S is inductively defined as follows:

$$\begin{aligned} CD(\neg\varphi) &= \text{negate}(CD(\varphi)) \\ CD(\varphi \wedge \psi) &= (\text{conjunctions}(CD(\varphi)) ++ \text{conjunctions}(CD(\psi)), \emptyset) \\ CD(\varphi \vee \psi) &= (\emptyset, \text{disjunctions}(CD(\varphi)) ++ \text{disjunctions}(CD(\psi))) \\ CD(\varphi \Rightarrow \psi) &= (\emptyset, \text{disjunctions}(\text{negate}(CD(\varphi))) ++ \text{disjunctions}(CD(\psi))) \\ CD(\varphi) &= (\varphi, \varphi) \text{ otherwise} \end{aligned}$$

with

$$\begin{aligned} \text{negate}([\varphi_1, \dots, \varphi_m], [\psi_1, \dots, \psi_n]) &= ([\neg\psi_1, \dots, \neg\psi_n], [\neg\varphi_1, \dots, \neg\varphi_m]) \\ \text{conjunctions}((V, W)) &= \begin{cases} V & \text{if } V \neq \emptyset \\ \bigvee_{w \in W} w & \text{otherwise} \end{cases} \\ \text{disjunctions}((V, W)) &= \begin{cases} W & \text{if } W \neq \emptyset \\ \bigwedge_{v \in V} v & \text{otherwise} \end{cases} \end{aligned}$$

N.B. This has not been implemented yet.